# Evaluation of structural and temporal properties of ego networks for data availability in DOSNs

Andrea De Salve · Barbara Guidi · Laura Ricci

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Abstract The diffusion of Online Social Networks (OSNs) OSNs, such as Facebook, Twitter, or Google+, are based has influenced the way of how people interact with each other. OSNs suffer of several drawbacks, one of the most important one is the problem of privacy disclosures. Distributed Online Social Networks (DOSNs) have been proposed as a valid alternative solution. DOSNs are Online Social Networks implemented in a distributed way, such as a P2P system or opportunistic networks. The decentralization of the control has issued several challenges, one of the main ones is guaranteeing data availability without relying on a central server. To define users' data allocation strategies, the knowledge of both structural and temporal characteristics of ego network are required. Unfortunately, the lack of real datasets limits the research in this field. The goal of this paper is the study of the behaviour of users in a real social network in order to define proper strategies to allocate the users' data on the DOSN nodes. In particular, we present an analysis of the temporal affinity and the structure of communities and their evolution over the time by using a real Facebook dataset.

Keywords P2P social networks, DOSN, data availability, temporal affinity, community detection

# **1** Introduction

Online Social Networks (OSNs) have hardly increased their popularity, becoming today the most used web platforms. During the years, they have created a new way of interaction/communication among people and they connect millions users that produce a significant amount of social content. The most part of current

Department of Computer Science, University of Pisa E-mail: desalve@unipi.it, guidi@unipi.it, laura.ricci@unipi.it

on centralized architectures which have intrinsic drawbacks including scalability and privacy [6]. Privacy is one of the most important issues regarding the usage of centralized OSNs because contents generated by users are stored and kept available by the centralized service provides. As a result, centralized OSNs have become the main channel of privacy disclosure. These drawbacks have led researchers to investigate alternative solutions, such as distributed approaches. A Distributed Online Social Network (DOSN) [6] is an online social network implemented on a distributed platform, such as a network of trusted servers, P2P systems or an opportunistic network. P2P paradigm has been largely used to implement DOSNs platforms where users' devices share the tasks needed to provide the OSN service. As a result, decentralization of the OSN allows to limit the privacy problem with respect the centralized service providers, but it introduces several new challenges such as data availability, information diffusion, or load balancing.

The extensive use of Smartphones and Tablets has affected the usage of OSNs and, nowadays, the most part of users connect to the OSNs by using their mobile devices. However, the possibility of using a wide variety of wireless radio access technologies (Bluetooth and WiFi) allows users' devices to be directly interconnected between each other via wireless links. The resulting networks established between mobile nodes (also referred as Mobile Ad hoc Network or MANET) allow their users to explicitly connect to the devices of the other users, in order to share information and contents. These features make P2P networks very dynamic over time and nodes participating in the system have to provide a data availability mechanism facilitating the seamless delivery of the contents.

Indeed, one of the major problems in a DOSN is to guarantee the availability of data, because of the high users' churn, i.e. the user may autonomously disconnect from the DOSN service and data stored on their is no more available and accessible. As a result, the behaviour of participating users dramatically affects the data availability and it can lead to data becoming unavailable or lost. Replication is one of the most popular approaches to manage this problem and it consists of replicating the user's contents on different nodes of the DOSN. Replication improves data availability, but, at the same time, it introduces the problem of consistency and the problem of minimizing the number of replicas.

For these reasons, the study of good policies for the selection of nodes where to store replica of contents has recently gained momentum. The goal of these politics is to dynamically allocate data such that the availability of the social data is maximized.

Existing DOSNs ensure data availability by different approaches. One solution is to encrypt social data and to store it on a Distributed Hash Table (DHT) [19] (such as [12, 1]). The other solution is to exploit the social relations to store data on the devices of friend nodes. Several solutions consider all friends, a set of trusted friends, or random subset of the friends [13, 8, 5, 20].

In this case, information on both the time spent online by users and the ability of the users to serve the data are also considered during replicas selection. Besides these, the study of data availability problem can benefit from information about interconnections between users of the DOSNs, i.e. communities resulting from social relations. Indeed, a very common trait of OSNs is to expose communities structures where groups of users (or clusters) are highly connected intracommunity and weakly connected inter-community.

These communities may reflect the different environments where people met (such as workplace, school, hometown, etc.) or the different features that aggregate them (such as interests, music, religion, etc.). Having a better understanding of such community structure occurring in DOSNs can help scientists in managing the data availability problem. Community structure may depends exclusively on the topological structure of the ego network, i.e. the friendship relations between the alters of the ego. Furthermore, it is interesting to analyze both the static structures of the communities and how they evolve when users connect/disconnect from the social network.

In addition, the obtained communities can be overlapped, another attractive property which models the real world assumption that a user may belong to more than one community at same time and can act as a bridge allowing communication of social information from one community to the other. From the perspective of data availability problem, these properties of OSNs users remain unexplored in the context of DOSNs.

The main goal of this paper is to present a twofold study by using a real dataset of Facebook. Firstly, we extend the analysis proposed in our previous work [7] regarding the temporal behaviour of users by focusing on the ego networks of individual users. In particular, we evaluate the extent to which friends of a user have the same temporal behavior, i.e. they are connected to the OSN in the same interval of time. Understanding how users navigate and interact when they connect to social networking sites creates opportunities for improvements in term of design how to manage the availability problem.

As a second step, we introduce a novel study concerning communities belonging to the ego networks of the users. First of all, we study static communities, i.e. communities which can be detected by considering the topological structure of the ego network. In addition, we investigate a dynamic scenario in which we detect the communities arising when we consider only the users which are online at a certain time step. Next, we perform a detailed analysis aimed at studying properties and temporal evolution of communities in each ego network, giving particular focus to a comparison of the dynamic scenario with respect to the static scenario. More in detail, we evaluate how communities structure in ego networks evolve when availability status of the users (online/offline) is considered. The analysis of the dynamic scenario confirm the presence of tightly connected group of nodes, reflecting different communities. In addition, dynamic communities became more sparse and poorly connected compared to the dynamic scenario.

At the best of our knowledge, this is the first contribution of this kind in the DOSNs field and, more in general, this work represents the first contribution about the study of both temporal ego networks and temporal communities detected by considering the online behaviour of users. Furthermore, the analysis detects the users belonging to more than one community that may act as "hubs", able to serve accesses to the replica they store for many users of the ego network. We consider the social communities defined by the social relationships. Thanks to communities, we can be able to guarantee a high level of data availability by replicating social data on nodes belonging these communities.

The rest of the paper is organized as follow. Section 2 describes the related work. Section 3 describes the scenario in which we focus and introduces the problem of data availability in DOSNs. Section 4 describes our twofold work on temporal affinity and community detection. Section 6 investigates the results of our analysis, which are obtained by exploiting a real OSN data set derived from Facebook. Finally, Section 7 draws the main conclusions.

#### 2 Related Work

Since our study is targeted to the DOSNs scenario, we introduce the reader to research fields related to current DOSNs and to the current proposals of techniques to guarantee data availability in DOSNs.

## 2.1 Distributed Online Social Networks

A Distributed Online Social Network (DOSN) [6] is an Online Social Network implemented on a distributed information management platform. Authors in [6] propose a reference architecture, which could consist of six layers where the lower layer is the physical communication network, which can be the Internet or another physical communication network and the distributed overlay which manages all the functionalities of the system is usually a P2P network. Specifically, this layer provides services for looking up resources, routing, and retrieving information reliably and effectively among nodes in the overlay.

Current proposed approaches can be divided into two major categories: Federation of servers and OSN over P2P networks.

OSNs built on top of a federation of servers require that social network providers agree upon standards of operations in a collective fashion. An example of these kind of DOSNs is Diaspora [8]. Instead, DOSNs built on top a P2P network represent the most studied category. A common choice is to use a DHT to support the DOSN. LifeSocial [12] is a P2P OSN focused on the privacy issue, where user information is stored by exploiting a DHT and it presents an approach where all OSN functionalities are realized by plug-ins.

PeerSon [1] is a distributed infrastructure for social networks whose focus is related to security and privacy concerns. It proposes a two-tier architecture where the first tier is a Distributed Hash Table (DHT) and the second tier consists of the nodes representing users. The idea is to use the DHT to find the necessary information for users connecting directly to the target nodes. This approach comes without a replication scheme and stores offline messages on the DHT (OpenDHT in the prototype implementation). All users' content is encrypted.

Safebook [5] addresses privacy in OSNs by using a three-tiers architecture where data are stored in a particular social overlay named "Matryoshka". Matryoshkas are concentric rings of peers around each users peer that provide trusted data storage and obscure communication through indirection.

#### 2.2 Data Availability in DOSNs

One of the main challenges in decentralization comes from guaranteeing availability of the data when the owner of the data is not online. When we talk about data availability, we talk about an important challenge in distributed systems. Without a central point of storage, data are distributed among all nodes in the network. When a node leaves the network, its data should remain available in the network. In the context of DOSNs, we consider availability of user's data, called content, that can be seen as the digital representation of a user, which is stored on a computing device and can be transmitted from one device to another.

The main technique used to resolve the problem of the content availability in DOSNs is the Replication [35]. Replication is a well-known technique in a distributed system and it is a technique based on the storage of the same data on different storage devices.

In DOSNs, probably due to the lack of real data, only few proposals have addressed this issue. Super-Nova [20] is an architecture for a DOSN that solves the availability issue by relying on super-peers that provide highly available storage. Other specific solutions are presented in [15], where authors propose a replication strategy based on storing the replicas of users profiles only on a set of trusted proxies.

A current trend to manage the data availability problem is the usage of social relationships to store social data. DiDuSoNet [13] is a P2P DOSN focused on the data availability problem which uses, as SafeBook, a particular overlay based on trusted connection to store data.

The approach proposed by Koll et al. [14] is SOUP, a scalable, robust and secure system for storing data among heterogeneous nodes with very high availability, which considers the significantly differing characteristics of OSN users, such as device and the online time patterns.

# **3** Scenario and Motivations

We consider a scenario where the support of the DOSN is P2P network which exploits an overlay derived from the social relations existing between the users. In particular, we refer to a specific social network model defined as *ego network* and explained in detail in [13]. This model is well suited for this purpose, representing a local view of the social graph relative to a user, where we consider only his/her friends and their interconnections. To be more detailed, the ego network [10] is the network constituted by a user (*the ego*), its direct friends (*the alters*) and the social ties occurring between them.

In our previous works [3], we have studied replication strategies based on both friendships and temporal information. Taking into account both these information, we are able to maximize the data availability by choosing good replica nodes. In this work, we are interested in studying the temporal affinity between Facebook users and their friends by considering the ego network structure. Furthermore, we study how communities of nodes evolve if we consider the presence status (online/offline) of the users. Indeed, the presence of densely connected groups of nodes can be exploited to increase the level of data availability. For these reasons, we first study the static ego network to evaluate some characteristics of the individuated communities. The second step is evaluate the characteristics of these communities over time by using the presence status of the users.

# 4 Ego network analysis: temporal affinities and static communities structure

In this section, we introduce our studies concerning the users' temporal affinity and the detection of static communities.

# 4.1 Temporal Affinity

With the term *Temporal Affinity* we refer to the phenomena of users having the same temporal behaviour, i.e. the probability that they are online in the same interval of time. In more detail, considering the ego network of an ego E, we study how similar is the temporal behaviour of the alters in his/her ego network with respect to that of E. The availability of each user is represented by an availability vector of fixed size.

To analyse the temporal affinity, we use a specific *presence array* of 2001 entries (one for each temporal slot in our dataset) for each couple ego-alter. Each slot refers to 8 minutes and it contains a value: 1 if the user is online in this slot, 0 otherwise. Since our goal is to understand how the temporal behaviour of users can affect the data availability in DOSNs and considering that there are particular day periods in which users tends to be offline (e.g. during the night), we propose two different metrics for the Temporal Affinity:

- Daily Temporal Affinity (DTA), in which we analyze the temporal affinity during all the day. In detail, we exploit all the 2001 temporal slots;
- Nighttime Temporal Affinity (NTA), which considers a subset of nighttime slots.

We evaluate the temporal affinity using the cosine similarity metrics [2]. Formally, let A and B the availability vector of two users, the cosine similarity is computed as shown in (1):

CosineSimilarity
$$(A, B) = \frac{A \cdot B}{||A|| \cdot ||B||}$$
 (1)

The resulting similarity ranges from 1 meaning perfect correlation, to 0, usually indicating no correlation between the ego and the alter.

To evaluate the NTA, we need to define when an ego can be considered active during the night (i.e from 12:00 midnight to 06:00 a.m.) of a day *i*. An ego is *nighttime* during the night of the day *i* if and only if it has been online for at least 15 temporal slots. We define an *knighttime ego* as an ego which has been nighttime for *k* nights and a *nighttime alter* as an alter which is online for the 95% in the same nighttime slots of a *k*-*nighttime ego*. The *nighttime affinity coefficient* is defined as the average number of nighttime alters of a *k*-*nighttime ego*.

#### 4.2 Community Detection

Communities structure is considered to be a significant property of social networks. Often complex networks, such as OSNs, are characterized by densely connected groups of nodes which are sparely connected to other nodes, named communities. A user usually has connections to several social groups like family, friends, and colleagues. Further, the number of communities an individual can belong to is actually unlimited because a person can simultaneously take part in as many groups as he/she wishes [22].

Numerous techniques have been developed for both efficient and effective community detection. One of the most complete survey is given by Fortunato [11], which describes all the most popular techniques for community detection. One of the most common ones is based on the Label Propagation algorithm (LP), where social relationships between users of the DOSNs are exploited to identify communities of users in a distributed and decentralized way. To discover communities in the ego networks, we decide to use the Label Propagation (LP) [17], implemented by DEMON [4]. Label Propagation, according to the original work, detects communities by spreading labels through the edges of a graph and by labeling a node with the label paired with most of the neighbours of the node. DEMON implements the classical label propagation algorithm, as explained by [17], but, Unlike the original algorithm, DEMON is able to identify possible overlaps between communities. Furthermore, DEMON is suitable for a distributed implementation and it can be easily adapted to any kind of network, dense or sparse.

To evaluate the overlapping between communities, we define a new index, the *k*-overlapping index (KOI). Taking into account the definition of an ego network as given in Section 3, the index is used to understand how many alters of a generic ego e belong to at least k communities. In detail, the index is defined as follow: consider  $C_S$ , the set of communities detected in a generic ego network EN, for each alter node  $a \in EN$ , we compute the number of communities in  $C_S a$  belongs to. We detect alters that appear in at least k communities and compare this value with the total number of nodes in the ego network. A formal definition of the KOI is given by (2).

$$KOI(EN,k) = \frac{|V^k|}{|V|} \tag{2}$$

where V is the set of alters nodes in EN and  $V^k$  is the set of alters in EN which belong to, at least, k communities.

On the basis of the KOI index, we can detect alters which belongs to a significant set of communities and are, therefore, suitable to host and spread replicated data. On the other way round, the allocation strategy must take into account that since each user device has a limited amount of memory, especially if we consider mobile nodes, allocation can cause a huge amount of load on a single node characterized by a high KOI index because it could be chosen by a large set of friends to replicate data.

# 5 Dynamic evaluation of ego networks and communities

Our consideration about community detection, presented in the previous section, can be applied, in part or as a whole, to a dynamic scenario. An evaluation of communities in a static scenario can be a very good starting point, but it is not enough when we consider the behaviour of the users in a DOSN. Of course, the temporal affinity between users just has to consider the presence status of the users, which evolves during the time. Instead, an important characteristic that has to be considered in a distributed environment is the time evolution of the communities in term of their composition. For this reason, we have evaluated how the dynamic structure of the communities can affect data availability, by



Fig. 1 Snapshot of the communities structure and availability status (online/offline) of the ego network of user E.

assessing how their change over time in term of number and size.

As explained in [16], the structure of most complex networks is dynamic in nature because of the addition and/or the removal of links and nodes. This has an impact on the communities which change over time. Consider a dynamic scenario as a distributed system in which nodes are represented by participating users, which use different devices but that are identify in the graph as one single node and, links are represented by the connection between users which can be more or less stable depend on the network infrastructure used (wireless network, etc..). In this scenario the dynamic nature of the network is really hard to manage because the presence status of a user drastically changes and connections/disconnections within a small time interval could be very huge. In such a case, it is important to detect online nodes that form a community to facilitate data availability and information diffusion. Indeed, dynamic, densely connected communities of nodes can be used to provide contents of one user to all the members of the community.

Consider the example of an ego network of a node and the communities detected in it, as shown in Figure 5. The different communities detected by the algorithm are shown with different colours. The communities are computed by considering the topology of the ego network, i.e. nodes belong to different communities, according to their social relationships with the friends of the ego.

Now suppose that the ego voluntary decides to leave the network at the time t. To choose the nodes where data of the ego has to be mapped, it is necessary to consider the snapshot (at time t) of all online nodes and to dynamically compute the communities on these nodes. As shown by Figure 5, communities structure may change frequently over time due to the disconnection/connection of users. Indeed, static communities can be considered as an optimal situation which occurs only when all the nodes of the ego networks as online. In contrast, communities arising when availability status of users is considered may appear very different because only part of the users of the communities are available.

To ensure the availability of data, user's contents must be replicated on different nodes (among those online), according to the friendships in the ego netwokrs. However, friendship relations give few information about the quality of the replica nodes. Instead, communities information provides a complete view of the structural properties of the ego networks and can be used to place replica of the contents on nodes located in different communities.

For this reason, the evolution of the presence status of the nodes of the ego network during the offline interval of the ego node must be monitored. We study the dynamic evolution of ego networks by considering both the evolution of the presence status of the nodes during the time period in which an ego is offline.

# 6 Evaluation

In this section, we describe our evaluation concerning the temporal affinity and the communities belonging to an ego network.

We have used a real dataset gathered by a Facebook application, named *SocialCircles!*, deeply described in [7]. At the best of our knowledge, our dataset is the only one which contains structural and temporal information about users.

We sampled 337 registered egos and their friends every 8 minutes for 10 consecutive days (from Tuesday 3 June 2014 to Friday 13 June). Using this methodology we were able to access the temporal status of 308 registered users and of their friends (for a total of 95.578 users). For the purpose of clarity, we will refer to ego nodes to indicate these 308 users. We consider the availability trace of each user to determine the start of a session and its termination. More specifically, time starts at the beginning of time  $s_0$  and it is segmented in a subsequent time periods (each of 8 minutes), that we refer as *time slots* in the follow.

# 6.1 Temporal Affinity

The first analysis concerns the evaluation of the Daily Temporal Affinity, referenced as DTA. As explained in [7], a user can be online, offline or in a idle state.

We decide to consider the online state, and to do not distinguish the idle and offline state into the presence



Fig. 2 Daily temporal affinity

array so that both of them are represented by the value 0.

About the 80% of the couple ego-alter has a low similarity, as depicted in Figure 2. This low values can be influenced from the online behavior of each ego. A little set of couples (less than 5%) show medium/high values of similarity. These results confirm our previous analysis on Dunbar-based ego networks [7] where users present a high similarity only with trusted nodes. In fact, extending the temporal affinity analysis to the whole ego networks, the result is similar. For sake of clarity, Dunbar [9, 18, 21] shows that the number of friends a user can maintain stable social connections with is approximately 150 (the *Dunbar's Number*). The stability of a connection can be defined as a function of the tie strength of the relation and it quantifies the strength of the relationship between two users.

The second step concerns the analysis of the ego networks to find the Nighttime Temporal Affinity, indicated with NTA through the *Temporal Affinity index* introduced in Section 4. For the evaluation, we consider both the online and the idle state.

Some periods of the day are critical to guarantee data availability in distributed environments, such as the night period. When an ego has a high NTA with a set of distinct alters, it means that the tie strength between the ego and these alters is strong.

For the evaluation, we use two indexes: fixed indexes, when the night is divided into 38 slots from about the 03.00 A.M. to the 08.00 A.M. and variable indexes, when the night is defined by the slots containing less than 10000 users. The value k indicates the number of nights needed to consider an ego as a nighttime ego.

Figure 3 shows the CDF of the NTA of users who appear to be online for at least k = 1 night and k = 4 consecutive nights. About 90% of the nighttime users have at most 10 alters who behave in the same way, i.e.



Fig. 3 NTA with fixed indexes, by varying the number of nights k

alters who apper to be online during the nigh times. In general, nighttime users and alters does not changes their behavior over time, distributions of the NTA are very similar, regardless of whether the users are online for k = 1 or k = 4 consecutive nights.

We focued on variable index that consider extended periods of the day and computed the cosine similarity between availability vector of ego and alters who appear to be online during the night. Figure 4 show the CDF on the cosine similarity as well as the percentage of alters in each ego network. We observe that only few alters have availability pattern of the nighttime users. Analyzing the NTA we notice that less than 10% of egos can be considered nighttime. This value is drastically reduced when the constraint of the number of nights needed to be nighttime increases. For example, if we consider an ego network of 400 nodes, the number of alters which contribute to the NTA is less than 40, considering the state online and idle and this number decreases when we consider only the online state (less than 10). This results tend to confirm the Dunbar's circles [9] conjecture, showing that the number of nodes which we interact with is very low if compared with the total number of alters.

# 6.2 Detection of Static Communities

Our goal is to identify, in each ego network, a group of users which are densely connected among them. We have considered the ego network EN of each of the 337 registered users and we have executed DEMON on EN to compute the number of communities in EN. As expected, the dataset has revealed the typical structure of the social network, including groups of nodes which are strongly interconnected between them (high cluster coefficient).

Figure 5 depicts the CDF of the community size and of the number of communities. About the 80% of communities has less than 250 nodes. The ego networks show a complex structure and about 80% of them has a number of communities less than 15 (Figure 5(b)). This information is important when we consider the goal of our analysis. In fact, it permits us to estimate the number of replicas which can be allocated for the ego's data. Another interesting analysis evaluates the overlapping of communities. We exploit the *k*-overlapping index defined in Section 4 to evaluate, for each ego, the number of communities an alter belongs to. To choose the value k, we consider the average number of communities an ego node could have, as shown in Fig. 5(b), varying kfrom 2 to 5.

Figure 6 shows the evaluation of the overlapping index in communities. Communities present a considerable overlap, also when we consider increasing values of k. This means that ego networks have a set of nodes which belong to more than 2 communities, these nodes represent a bridge between different communities and are central in the ego network. In detail, when we consider the two bounds of k (Figure 6(a) and Figure 6(d)), we can clearly notice low values of the KOI index for k equal to 5, but we note that the increase of k is not proportional to the decrease of the index.

The properties we have detected for the communities highlight a guideline for the definition of a strategy for allocating the users' data replica in the ego network. A suitable node to store replicas of the user's data may be chosen among the set of nodes which belong to more than one community. A user paired with one of these nodes has direct friendship connections with many other nodes of the ego network and may guaran-



Fig. 4 Variable indexes with k equals to 2



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Fig. 5 Community Detection: an analysis

tee a good coverage of the ego network. On the other hand, a clever allocation has to take into account also the problem of load balancing, because nodes belonging to a set of communities could be affected by a huge load.



6.3 Dynamic evolution of communities

We have used the knowledge about the presence status of each ego node to build 2001 snapshots, each one related to a time slot. These snapshots are used to study the properties of the ego networks during the time and the evolution of the communities structure.

Table 1 gives some details concerning the comparison between dynamic ego networks and the corresponding static view of the same ego networks. Results suggest that the size in term of both nodes and links of the dynamic snapshot of the ego network is about 10%of the static view. In general, about 8% of all nodes in the ego network are online simultaneously which accounts for about 6% of the total friend relationships. This is motivated by considering the huge dynamism of ego networks. An important result is highlighted by the degree centrality which shows the presence of central nodes. The values are higher than the static ego networks because of the number of nodes are lower and nodes seems to be more central. Moreover, the ego networks do not result densely connected by considering that the average clustering coefficient which is 0.333. An interesting property is disclosed by the degree centrality. In fact, the average value of the degree centrality is 0.4, which means that there are nodes that are central and can be consider as hubs inside the ego network.

Figures 7(a) and 7(b) show in more details the evolution of, respectively, the number of nodes and of edges during each time slot. The plots clearly indicate the presence of a temporal pattern where most of the users are connected during daylight hours. In Figure 7(c), we evaluate the *local degree* that represents the average number of common neighbours between alters and the



Fig. 6 Community Overlapping through the k-overlapping index

	Dynamic			Static				
	Max	Mix	Mean	StdDev	Max	Min	Mean	StdDev
Num. Nodes	379.0	1.0	38.076	31.833	2965	23	486.89	361.60
Num. Edges	14513.0	0.0	663.007	1213.0118	256968	62	10930.14	22218.14
Degree Centrality	0.9166667	0.0	0.400	0.365148	0.88	0.0003427	0.0656477	0.07920130
Clustering Coeff.	1.0	0.0	0.333	0.5773	0.88	0.42	0.63	0.078

 Table 1 Comparison between static and dynamic properties of ego networks.

ego. Ego networks appears to be more clustered, online alters are connected on average to at least 2 commons online alters.

Let us now introduce the *Transition Interval* as the time interval when a generic ego node e decide to go offline and remain offline. For the sake of clarity, this is the interval starting from the last time slot when the ego node e is online to the last time slot when the ego node results offline. For example, if we consider a small presence array of 10 time slots 1110000111, the *Transition Interval* is the sequence 10000. In the following analysis, we consider the global view of the ego networks and what happens in the ego networks by focusing only on all the Transition Intervals of the ego.

Figure 8 shows the behaviour of the detected communities. Figures 8(a) and 8(b) show, respectively, the number of communities and their size evolving over time, not considering the presence status of the ego. Results clearly indicate that we can expect that the availability status of the ego (online or offline) does not strongly affect the communities structure of his/her ego network. In particular, ego networks expose community structures even when the presence status of users is considered. In general, the average number of communities is quite big (400 communities distributed among the different ego networks) while their size ranges between 10 and 20 users. In particular, the number of communities arising when users is offline is quite big (on average 250 communities) and they are composed by a num-



Fig. 7 Evaluation of the dynamic ego networks

ber of online nodes that does not exceed 20 users (see Figure 8(c) and 7(a)). In some cases (generally during the evening and in the early afternoon) communities present higher number of online members.

Table 2 shows some statistics of the dynamic communities as a function of the size of their ego network and compare them with the corresponding static ego network. For the static case, results suggest the presence of a significant group of communities in each ego network. We can observe that big ego networks (with a number of friends greater then 500) have a low number of communities with high size, while small ego networks (with a number of alters between 10 and 500) have a higher number of communities with a low dimension. The size of the communities is correlated to the number of alters and nodes of each communities are highly interconnected to each other. The scenario is completely different when we consider the dynamic case and many important characteristics are visible by comparing the two cases. In the dynamic scenario, the larger is the ego network size, the larger is the average number of communities. This means that large ego networks tend to

be more fragmented if we consider the time evolution. This depends also by the fact that big ego networks tend to have big communities as shown in the static scenario. Due to the users' churn, these communities tend to become fragmented over the time slots. Another important characteristic is highlighted by the average size of the communities which increases as a function of the ego network size. However, the gap is reduced and the difference between small ego networks and big one is of a factor of 4.

# 7 Conclusion and Future Works

This paper has analyzed the temporal behaviour of users and the communities of a real OSN to understand how these issues can affect the data availability in DOSN. We have noticed that not all the egos have a similar behaviour and that only a subset of nodes in their ego networks are good candidates to store replicas of profiles, by confirming the results obtained in [7] The analysis has reported interesting results when the temporal affinity has been evaluated during the night.



(a) Number of Community by considering all the ego (b) Size Community analysis by considering all the ego nodes



(c) Number of Community by considering Transition In- (d) Community Size by considering Transition Intervals

Fig. 8 Evaluation of the number and size of communities discovered in each ego networks: regardless of the online status of ego (Fig. 8(a) and 8(b)) and when ego is offline (Fig. 8(c) and 8(d)).

**Table 2** Average number of communities and average size by comparing ego networks with different size for both the Staticand Dynamic mode

	Dynar	nic	Static		
EN Size	Avg. Num. Comm.	Avg. Size	Avg. Num. Comm.	Avg. Size	
$0 \sim 199$	7	6.292697247706	4.4	44.01	
$200 \sim 299$	12	6.325449473	7.8	65.1	
$300 \sim 399$	20.7	8.0439574713	9.4	87.71	
$400 \sim 499$	31.13	8.92754612	11.4	103.02	
$500 \sim 699$	38.34	12.103263638	11.4	141.03	
$700 \sim 899$	87.40	13.8701502	15	175.51	
$900 \sim 2999$	132.76	27.80695214	16.4	342.75	

Furthermore, we have investigated the static communities in the ego networks and shown that they are heterogeneous in term of structure and size. We found that ego networks are characterized by a low number of communities, which does not depend from the ego network size and that there is an high level of overlapping of these communities.

The analysis has been extended with the evaluation of the evolution of communities structure in ego networks. In particular, we have investigated the behaviour of the communities over time, when the ego is offline. We have detected that ego networks expose communities structures, even if online/offline status of users is considered. More in detail, our work showed there is an evolution of the dynamic communities structure in ego networks, which became smaller and poorly connected, with a low clustering coefficient compared to the static communities. However, analysis of dynamic ego networks confirmed the presence of tightly connected groups of nodes, reflecting different aspects of the ego life. These results can be used to support a proper data allocation strategy, for example by using the number of overlapping communities as parameter of the selection strategy. We plan to extend our work in several directions, from a more systematic study of the communities for a better characterization of their features, to the definition of a good strategy to select the nodes to store data replica.

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